



## An Algorithm for Secure Key Distribution in Manets

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**Abstract**— a mobile ad hoc network (MANET) is self-organizing, dynamic topology network, which is formed by a collection of mobile nodes through radio links. MANETs enable wireless communication between mobile devices without relying on a fixed infrastructure. Hence, routing in dynamic network is a new challenge. We used various routing algorithms for smooth exchange of information between mobile nodes. Generally we imply security on all the nodes of the network. But this causes wastage of time and cost. This paper proposes that first we have to find out the shortest path and then imply the security in multicast routing. The routing is done with the help of GDH. We use the Key obtained by GDH to transfer data from source to destination. But there can be a malicious node that can change the key so that it can change the data sent to the destination node. So we have to find out these intruders (malicious nodes) that i have proposed in this paper. This proposed scheme will improve the performance of the network such as delay and packet delivery ratio than traditional routing algorithms.

**Keywords**— Mobile ad hoc network, group key management, multicast, GDH, Cryptography.

### I. INTRODUCTION

Mobile ad hoc network (MANET) consists of nodes which are connected by wireless links, where each node communicates with other nodes directly or indirectly through intermediate nodes. Thus, all nodes in a MANET basically function as mobile routers participating in some routing protocol required for deciding and maintaining the routes [1]. Routing in MANETs is challenging since there is no central coordinator that manage routing decisions. Routing is one of the key issues in MANETs due to their highly dynamic and distributed nature. Numerous ad hoc routing algorithms exist to allow networking under various conditions. They can be separated into three groups, proactive, reactive and hybrid algorithms [2]. In proactive routing algorithms maintain continuously updated state of the network and the existing routes; however, in some cases it may generate an unnecessary overhead to maintain the routing tables and then may be better to create routes only on demand, the case of reactive routing algorithms. In Reactive routing algorithms require time-consuming route creations that may delay the actual transmission of the data when sources have no path towards their destination and then, in this case may be better to use a proactive routing algorithm. In hybrid protocols try to profit the advantages of both reactive and proactive protocols and combine their basic properties into one. These protocols have the potential to provide higher scalability than pure reactive or proactive protocols thanks to the collaboration between nodes with close proximity to work together and therefore reduce the route discovery overhead multiple routing protocols have been developed for MANETs. In proactive\ protocols, every node maintains the network topology information in the form of routing tables by periodically exchanging routing information. Routing\ information is generally flooded in the whole network. Whenever a node requires a path to a destination, it runs an appropriate path finding algorithm on the topology information it maintains. The destination sequenced distance vector routing (DSDV) protocol and wireless routing protocol (WRP) are some examples for the proactive protocols. Reactive protocols do not maintain the network topology information. They obtain the necessary path when it is required, by using a connection establishment process. Hence these protocols do not exchange routing information periodically. The dynamic source routing (DSR), Ad-hoc on-demand source routing (DSR), Ad-hoc on-demand distance vector routing (AODV) [7], and temporally ordered routing (TORA) algorithm are some examples for the protocols that belong to this category. In proactive routing algorithms maintain continuously updated state of the network and the existing routes; however, in some cases it may generate an unnecessary overhead to maintain the routing tables and then may be better to create routes only on demand, the case of reactive routing algorithms. In reactive routing algorithms require time-consuming route creations that may delay the actual transmission of the data when sources have no path towards their destination and then, in this case may be better to use a proactive routing algorithm. In hybrid protocols try to profit the advantages of both reactive and proactive protocols and combine their basic properties into one.

### II. Distributed Bellman-Ford

The DBF algorithm was developed originally to support routing in the ARPANET. A version of it is known as Routing Internet Protocol (RIP) [1] and is still being used today to support routing in some Internet domains. It is a table-driven routing protocol, that is, each router constantly maintains an up-to-date routing table with information on how to reach all possible destinations in the network. For each entry the next router to reach the destination and a metric to the destination are recorded. The metric can be hop distance, total delay, or cost of sending the message. Each node in the network begins by informing its neighbours about its distance to all other nodes.

The receiving nodes extract this information and modify their routing table if any route measure has changed. For instance, a different route may have been chosen as the best route or the metric to the destination may have been altered. The node uses the following formula to calculate the best route:

$$D(i, j) = \min [d(i, k) + D(k, j) ]$$

Where  $D(i, j)$  is the metric on the “shortest” path from node  $i$  to node  $j$ ,  $d(i, k)$  is the cost of traversing directly from node  $i$  to node  $k$ , and  $k$  is one of the neighbours of node  $i$ . After recomputing the metrics, nodes pass their own distance information to their neighbour nodes again. After a while, all nodes/routers in the network have a consistent routing table to all other nodes.

This protocol does not scale well to large networks due to a number of reasons. One is the so-called count-to-infinity problem. In unfavourable circumstances, it takes up to  $N$  iterations to detect the fact that a node is disconnected, where  $N$  is the number of nodes in the network [2]. Another problem is the increase of route update overhead with mobility. RIP uses time triggered (periodic, about a 30-s interval) and event-triggered (link changes or router failures) routing updates. Mobility can be expressed as rate of link changes and/or router failures. In a mobile network environment, event-triggered routing updates tend to outnumber time-triggered ones, leading to excessive overhead and inefficient usage of the limited wireless bandwidth.

### III. Diffie-Hellman Two-Party AGREEMENT (DH)

This basic protocol, proposed in a landmark paper [3], allows two nodes to build a common key. The principal of this protocol is simple: the two involved nodes,  $M_1$  and  $M_2$ , send one another a partial key to be used for the common key computation.  $M_1$  generates a random number  $r_1$  ( $1 \leq r_1 \leq p$ ), and sends  $\alpha^{r_1}$  to  $M_2$ , such that  $\alpha$  and  $p$  are constants known by each node. On the other hand,  $M_2$  generates a random number  $r_2$ , and sends  $\alpha^{r_2}$  to  $M_1$ . Thereby, each node could compute the common key, which is  $\alpha^{r_1 \times r_2}$ . This solution is based on discrete logarithmic arithmetic, and also relies on the agreement on the parameters  $\alpha$  and  $p$  between the two nodes. Although it is simple and limited to two nodes' common key establishment, this protocol was used to design more sophisticated protocols, as we will see later [5]. This protocol uses exponentiation to share a secret between two parties, Alice and Bob. The protocol involves an initiator, Alice, and a responder, Bob. We use the common notation  $A \rightarrow B: M$  to stand for “A sends message M to B”. Raising message M to the power of exponent X is denoted by  $(M)^X$ . There is a public term denoted by  $g$ , which will be the base of our exponentiations. We represent the product of exponents by using the symbol  $*$ . Nonces are represented by  $N_X$ , denoting a nonce created by Principal X. The protocol description is as follows.

1.  $A \rightarrow B: A$   
Alice sends her name to Bob.
2.  $A \rightarrow B: B$   
Alice sends Bob's name to Bob.
3.  $A \rightarrow B: g^{N_A}$   
Alice creates a new nonce  $N_A$  and sends  $g^{N_A}$  to Bob.
4.  $B \rightarrow A: B$   
Bob sends his name to Alice.
5.  $B \rightarrow A: A$   
Bob sends Alice's name to herself.
6.  $B \rightarrow A: g^{N_B}$   
Bob creates a new nonce  $N_B$  and sends  $g^{N_B}$  to Alice.

Intuitively, when Bob receives  $g^{N_A}$ , he raises it to the  $N_B$ , to obtain  $g^{N_A N_B} = g^{N_A * N_B}$ .

Likewise, when Alice receives  $g^{N_B}$ , she raises it to the  $N_A$ , to obtain  $g^{N_B N_A} = g^{N_B * N_A}$ .

And due to the commutativity of the symbol  $*$ , they know the equivalence  $g^{N_B * N_A} = g^{N_A * N_B}$ . An observer of the exchange who does not know  $N_A$  nor  $N_B$  cannot find  $g^{N_A * N_B}$ , and so Alice and Bob have computed a shared secret, i.e.,  $g^{N_B * N_A}$ . Of course, the attacker can always learn a term  $g^{(N_A * N_I)}$ , where  $N_I$  is a nonce created by the intruder, even by using a passive intruder model. The point is that he can also make believe to Alice that  $g^{(N_A * N_I)}$  is the shared key she is sharing with Bob. This is usually modelled by adding to the protocol a new message where Alice sends to Bob some secret, encrypted by  $g^{(N_A * N_I)}$ . Existence of an attack is expressed by saying that the attacker can obtain this secret. For the sake of simplicity and because we are focused in AC-theories, we omit this last part of the protocol and concentrate just in whether the intruder can learn  $X^{N_A}$  for some exponentiation X, where  $X^{N_A}$  is the key calculated by Alice.

In a rule-based representation of this protocol, parts of a received message whose make-up cannot be verified by a principal are represented by variables. That is, since nonces are known only to the principal who generated it, and retrieving the nonce would require the computation of a discrete logarithm, we say that Bob receives a variable X of a generic message sort instead of  $g^{N_A}$  and similarly for Alice. The symbol  $*$  is associative and commutative and satisfies the following additional property with respect to exponentiation:  $(X^Y)^Z = X^{(Y*Z)}$

The intruder abilities to create, manipulate, and delete messages according to the Dolev-Yao attackers capabilities [6] are described as follows, where we use the special symbol  $\underline{\mathcal{E}\Gamma}$  to represent that the intruder knows something, and I denotes the intruder's name:

$$\frac{M_1 \in \underline{\mathcal{E}\Gamma}, M_2 \in \underline{\mathcal{E}\Gamma}}{(M_1 * M_2) \in \underline{\mathcal{E}\Gamma}} \quad \frac{X \in \underline{\mathcal{E}\Gamma}, Y \in \underline{\mathcal{E}\Gamma}}{X^Y \in \underline{\mathcal{E}\Gamma}} \quad \frac{}{N_i \in \underline{\mathcal{E}\Gamma}}$$

The intruder also knows the names of all the principals and the base g.

If we ask ourselves whether the intruder can learn a message  $X^{N_A}$  for some variable X received by Alice (representing the nonce that Alice receives from Bob), the answer is yes for an infinite set of instances for X, e.g.,  $g^{N_i}, (g^{N_i})^{N_i'}, ((g^{N_i})^{N_i'})^{N_i''}$ , etc. If we take instantiation  $X \rightarrow g^{N_i}$ , the intruder can learn the message  $g^{(N_A * N_i)}$  by means of the following sequence of actions (only the three first steps are necessary but we need Alice to complete the protocol in order to believe she is sharing a shared key with Bob):

1.  $A \rightarrow B : A$   
Alice sends her name to Bob, but it is intercepted by the intruder.
2.  $A \rightarrow B : B$   
Alice sends Bob's name to Bob, but it is intercepted by the intruder.
3.  $A \rightarrow B : g^{N_A}$   
Alice creates a new nonce  $N_A$  and sends  $g^{N_A}$  to Bob, but it is intercepted by the intruder.
4.  $I \rightarrow A : B$   
The intruder sends Bob's name to Alice.
5.  $I \rightarrow A : A$   
The intruder sends Alice's name to Alice.
6.  $I \rightarrow A : g^{N_i}$

The intruder creates a new nonce  $N_i$  and sends  $g^{N_i}$  to Alice.

The intruder is able to learn the message  $g^{(N_A * N_i)}$  just by raising the intercepted message  $g^{N_A}$  to  $N_i$ . Note that the intruder does not need to know  $N_A$ , since he gets the desired effect thanks to the educational properties for exponentiation and product of exponents described above.

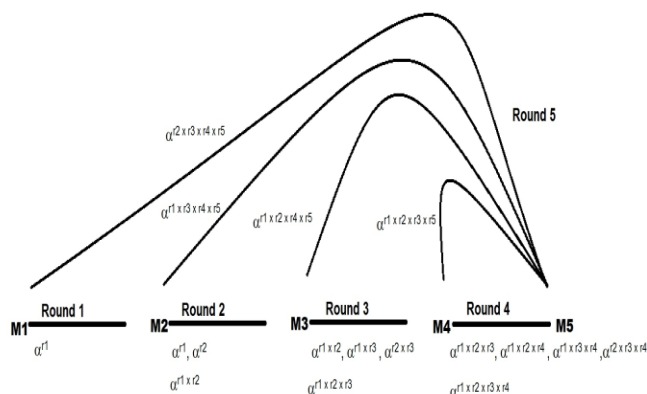
#### IV. General Diffie-Hellman (GDH)

Steiner *et al.* [4] proposed a n-party generalization of the basic two-party DH protocol (described before). The new protocol consists of  $n$  rounds, allowing  $n$  nodes to establish a common key. In the first  $n - 1$  rounds contributions are collected from each node. In the first round,  $M_1$  generates  $r_1$  and computes  $\alpha^{r_1}$ , which it sends to  $M_2$ . In the second step  $M_2$  generates  $r_2$ , computes  $\alpha^{r_2}$  and sends it to  $M_3$ , along with  $\alpha^{r_1}$  and  $\alpha^{r_1 * r_2}$ . This latter sends to  $M_4$  (after making the required computations) the third-round partial factors, i.e.,  $\alpha^{r_1 * r_2}$ ,  $\alpha^{r_1 * r_3}$ ,  $\alpha^{r_2 * r_3}$ , as well as the third-round partial key  $\alpha^{r_1 * r_2 * r_3}$ . This process continues for each  $M_i$  ( $i < n$ ). Upon the  $(n - 1)^{th}$  round, the collector node  $M_n$  receives the  $(n - 1)^{th}$  round partial factors, and the  $(n - 1)^{th}$  round partial key, then it generates its random number and computes the final key  $K$ . In the last round, node  $M_n$  sends each  $M_i$  the appropriate  $(n)^{th}$  round partial factor, i.e.

$$K = \alpha^{(\prod_{j=1}^n r_j) / r_j}$$

Consequently, each node uses its random number to compute the common key  $K$ . Note that partial factors are used to avoid sending the final resulted key during the last round. Also note that the  $(n - 1)^{th}$  round requires  $n - 1$  operations (sending the partial factor to each node), which makes the computational complexity of the solution  $O(2 \times (n - 1))$ .

Even though it uses a collector, this solution is contributory, since each node contributes to the key computation with the random number it generates. Nevertheless, the major drawback of this solution is the important overhead, due to the message size rising from round to round. This can also cause a problem with scalability.



Example of General Diffie-Hellman with five Node

Figure 1:Example of GDH

### V. Proposed Work:

Generally we imply security on all the nodes of the network. But this causes wastage of time and cost. This paper proposes that first we have to find out the shortest path and then imply the security in multicast routing. The routing is done with the help of GDH and the Encryption and Decryption of data is done with the help of a new Cryptographic technique that i have proposed in this paper. This proposed scheme will improve the performance of the network such as delay and packet delivery ratio than traditional routing algorithms:

- A network is randomly created
- Then the shortest distance from source to destination is found with the help of bellmen ford algorithm
- Apply GDH on each node starting from source to destination node on the shortest path calculated.
- Then Encryption and decryption will be done on the DATA from source to destination node.
- After that the Malicious node will be detected
- Then a new path will be found from source to destination.
- The new path will be based on nodes selected from authentication pool of nodes.
- The authentication pool consists of authenticated nodes.
- So that a secure path can be find out to transfer the data from source to destination nod.

### VI. Experiment Results And Analysis:

```
Enter The number of nodes
10
    0    1    2    3    4    5    6    7    8    9
0    0    94    9    28   61   97    4    7    17   40
1    94    0    32   78   29    8   49   60   48   42
2    9    32    0    3    62   52   70    5   70   42
3    28   78    3    0   10   40   47   33   46    1
4    61   29   62   10    0   11   90   56   90   26
5    97    8   52   40   11    0   38   58   44   27
6    4    49   70   47   90   38    0    2   96   60
7    7    60    5   33   56   58    2    0   37   56
8    17   48   70   46   90   44   96   37    0   86
9    40   42   42    1   26   27   60   56   86    0
```

```
Enter The Node with which u wanna communicate
8
```

```
0 -> 8 ->
```

```
Selecting Prime Numbers P and G
```

```
Secure Key: 506
```

```
*****
```

```
Checking the authenticity of network
```

```
Sending Data 506
```

```
Sending Encrypted Data though the network
```

```
Destination node received the message
```

```
Decrypting data
```

```
There is a middle man
```

```
*****
*****
```

```
Press 1 to continue
```

```
Press 0 to exit
```

```
1
```

```
Let there be a middle man in the Network
```

```
2
```

```
2
```

```
2
```

```
*****
```

```
Checking the authenticity of network
```

```
Sending Data 77
```

```
Sending Encrypted Data though the network
```

```
Destination node received the message
```

```
Decrypting data
```

```
There is a middle man
```

```
*****
*****
```

```
Press 1 to continue
```

```
Press 0 to exit
```

```
1
```

```
Authenticated Diffie Hellman
```

```
Enter The Node with which u wanna communicate
```

```
8
```

```
0 -> 8 ->
```

```
Authenticated Nodes are:
```

```
0    8    9    5    7
```

```
Authenticating Route
```

```
All nodes are authenticated
```

```
Selecting Prime Numbers P and G
```

```
Secure Key: 84
```

### VII. Conclusion:

Case I: If there is no intermediate node in between the source and the destination node. Then there is no way that the key shared between the two nodes gets intended and the data transfer is secure.

Case II: If there are some intermediate node in between the source and the destination node i.e. along the path between the source and the destination node and if some malicious node and if some malicious node intrudes the key and change it and then forwards it to the destined node waiting for the key to be shared, the destined node will not be able to know whether the key coming to be shared is either intruded or not and take it as a key that is not infected from any malicious node (intruder) .

Case III: Then the prime target is to detect the malicious node along the whole path and to make the path secure so that data can be transferred securely over the path without any intrusion.

Case IV: To avoid the malicious node we make an authentication pool of nodes in which we assume a threshold value of a node. If any node has that threshold value or more than that then it is an authenticated node otherwise not. So by doing this we can find the malicious node and can do the secure transfer of data.

### VIII. Future Work:-

We can do the clustering of the authenticated nodes so that we need not to check all of the nodes rather relevant cluster can be taken into account. By doing this we can save the time of choosing authenticated nodes and data transfer from source to destination.

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